DISTRIBUTED FRAMEWORK FOR FREQUENCY BASED THREAD POOL

Sahrish Imtiaz, Faisal Bahadur, Arif Iqbal Umar

# Department of Information Technology
Hazara University
Mansehra, Pakistan

Abstract—Concurrency is necessary in server side programming, and two basic ways to build concurrent programs i.e. single-threaded approach and Multithreading approach. Multithreading improves processor consumption in several ways, such as shared memory multiple processors (SMP), multiple threads can be executed on different processors to easily distribute the processing load. Thread pool is a multithreading architecture that is used in web and application servers to maximize performance and the dynamic optimization of thread pool is a challenge to retain server’s performance as the request rate on the server fluctuates. Frequency based dynamic thread pooling strategies face this challenge on vertical scaling where requests are processed by a single server. These strategies cannot be used in horizontal scaling for distributed applications hence unable to use the power of distributed servers. This paper presents a dynamic thread pool tuning strategy that is developed on a distributed framework to utilize the power of distributed servers named DFBOS that offers the advantages of substantial scalability and liveliness. The simulation results revealed the fact that DFBOS outperforms other frequency based thread pooling strategies in terms of throughput and response time.

Keywords—Concurrency, Concurrency Mechanisms, Multithreading Approaches, Restricted Frequency Based Optimization Strategy, Distributed Frequency Based Optimization Strategy

I. INTRODUCTION

Concurrency is the process of processing multiple jobs at same time in computer system. Typically two aspects are essential when dealing with concurrency. One is to control the external events occurring in random order and the second is to make sure that these events are making response in minimum needed time interval.

If every synchronous action changed severally, in a very actually similar approach, this is able to be comparatively simple: we've got an inclination to may simply produce distinct programs to run out all activities. The experiments of planning synchronous systems rise mostly as results of the contacts that occur among coincident actions. Once period of time activities move, some variety of coordination is required.

Concurrency control brings different benefits to the computer system i.e. exploiting multiple processors, improved resource consumption, more responsive programs and simple program design.

Multitasking

When two or more tasks are processed parallel at certain period of time. Multitasking is managed by Operating System. Computer System executes segments of multiple tasks in parallel, while the tasks share common resources such as CPU and RAM. By using the context switching mechanism, CPU execute a certain process for a specific interval of time, and after executing next process is loaded for processing.

Multithreading

When a CPU executes multiple threads at same time is called multithreading, which is supported by Operating System. By using Instruction Level Parallelism and Thread Level Parallelism, aim of multithreading is to increase the use of a single core. After 1990s Instruction Level Parallelism has become more popular in Multithreading execution. Threads are an approach of achieving a better quality of concurrency inside a process. Every thread of same process shares the same memory resources and other processing resources. Typically all threads are assigned a method to execute.

Multithreading Approaches

When Single CPU process Multiple Process or Threads is called Multithreading. Beside these statements,
synchronous blocking I/O is the archetypal way for running I/O. It is an archetypal style that is all around bolstered by abundant programming dialects. It likewise prompt a straight forward programming model, since all assignments vital for solicitation taking care of can be coded successively. Besides, it gives a basic mental deliberation by detaching solicitations and concealing simultaneousness. Genuine simultaneousness is accomplished by utilizing numerous Threads/forms in the meantime.

**Thread per Request Approach**

This approach manages every request from a client in a distinct thread of control. This model is convenient for servers that handle extensive duration requests such as database requests, from several clients. It is fewer helpful for short time requests due to the transparency of generating a new thread for every request. It can also consume a huge number of OS resources if numerous clients make requests at the same time.

**Thread Pool Approach**

In Thread pool framework, M quantities of Threads are made to do N quantities of errands. When all is said in done M≠N rather, the measure of Threads is tuned to the figuring assets accessible to handle undertakings in parallel (processors, centers, memory) while the quantity of errands relies upon upon the issue and may not be known forthright.

Purposes behind utilizing a Thread pool, as opposed to the undeniable option of producing one Thread for every undertaking, are to keep the time and memory overhead inborn in Thread creation, and to abstain from coming up short on assets, for example, open records or system associations (of which working frameworks designate a predetermined number to running projects).

A typical method for disseminating the assignments to Threads (booking the errands for execution) is by method for a synchronized line known as an undertaking line. The Threads in the pool take undertakings off the line, perform them, and after that arrival to the line for their next errand.

### II. RELATED WORK

This section gives the brief detail of various researches taken out on thread pool system. Till date a variety of thread pool models are anticipated by the researchers.

**“Object Interconnections by Douglas C. Schmidt & Steve Vinoski”**

This paper describes thread pool concurrency version and provide an explanation for the way to use it to develop multithreaded servers for a dispersed stock citation application. This paper describes how item-orientated strategies, C++, CORBA, and upper level abstractions like the Singleton pattern assist to make simpler programming and enhance extensibility. This paper explores another concurrency version: thread per consultation. This version is supported with the aid of some of CORBA implementations such as MT-Orbix and ORBeline. Having a preference of concurrency fashions can help developers meet the overall performance, functionality, and preservation necessities of their applications. The key to accomplishment, of direction, lies in very well knowledge the tradeoffs between distinctive fashions.

**“Design and Implementation of Multi-Threaded Object Request Broker by Winston Lo & Yue-Shan Chang”**

This paper explains CORBA 2.0 to design and placed into effect a multi-threaded entity request agent based totally on the Windows NT and critical TCP delivery protocol. This paper divides the ORB in 3 layers. These layers includes run time, records representation & communication layer. Run time layer is a set of energetic
petition interface exercises with which customer programs call a remote item. Inside the statistics illustration layer CORBA 2.0 is used, wherein commonplace statistics illustration is used & in the communication layer TCP/IP protocol is used, that's the standard of net. ORB uses the runtime library & server spirit to combine programs and item enforcement. The user interface is comparable to Orbix, as it compares performance with Orbix. It also implements an IDL compiler in ORB surroundings to interpret IDL definitions into C++ mapping. Then, it uses programs to calculate ORB overall performance via employing Schmitt’s approach to evaluate with IONA’s Orbix. From the results, it is clearly proven that the statistics marshalling, unmarshalling and greatest reminiscence allocation in statistics illustration layer in gadget is extra green than Orbix. Its miles as it adopts buffer pool reminiscence control to manage allocated memory.

Growing instructions of real time structures require end to give up support for diverse nice-of-provider (QoS) elements, including jitter, latency, bandwidth and dependability. Applications contain control and manipulate, mechanized manner manage, videoconferencing, massive-scale dispersed interactive simulation, and test beam facts acquisition. Those structures require support for Threadent QoS requirements. To fulfill this challenge, developers are turning to allotted object computing middleware, including the common item Request broker architecture, an object control organization (OMG) industry standard. In complicated actual-time structures, doc middleware is living between packages and the fundamental working systems, protocol stacks and hardware. CORBA enables lower the cycle time and struggle required to expand excellent systems with the aid of composing programs the use of reusable software program thing offerings in place of constructing them completely from scratch. Real Time CORBA specification includes functions to control CPU, network and reminiscence sources. The authors describe the key actual-Time CORBA features that they experience are the most pertinent to researchers and developers of disbursed real-time and implanted systems.

“Performance Study and Dynamic Optimization Design for Thread Pool Systems by Dongping Xu”
This paper provides a reason behind the set of efficiency metrics for quantitatively analysis of thread pool basic efficiency. For experiment, a thread pool device become constructed which presents a famous framework for thread pool studies. On the bases of this imitation environment, the general performance effect brought via the thread pool to amazing multithreaded applications. Furthermore, the correlations amongst inner specialties and the throughput turned into moreover studied the experimental consequences point out that the commonplace undertaking idle time has robust courting with the thread pool extraction. This paper projected and appraisals the idea of the usage of a heuristic method to resolve the quality thread pool period primarily based at the project common idle time. The simulation effects prove that dynamic optimization for thread pool period could be very powerful in eliminating the overhead and improving the overall efficiency.

“Dynamic Thread Count Adaptation for Multiple Services in SMP Environments by Takeshi Ogasawara”
This paper proposes a dynamic technique, thread count adaptation, which adjusts the thread counts that are owed to services for adapting to CPU requirement
variations in SMP environments. The target is to raise the utmost throughput accessible on a system that has numerous dynamic content services while meeting various service time criteria for these services in dynamic workloads. Confront is to noticeably progress response times for dynamic content on a busy well-tuned thread-pool-based system without prioritizing any specific services. The experiments reveal that a prototype using our approach on J2EE middleware quickly (around every 20 ms) adjusted the thread counts for the services and that it improved the average 90th-percentile response times by up to 27% (and 22% on average) for the SPECj AppServer2004 benchmark.

“A Dynamic Adjustment Mechanism with Heuristic for Thread Pool in Middleware by Ning-jiang Chen & Pan Lin”

The primary position of this paper is the dynamic alteration approach of thread pool pushed by means of heuristic actions (which incorporates response coefficient and jamming thread pointer) has been supplied. Those heuristic actions can replicate run time reputation of thread pool correctly; therefore the dimension of thread pool may be tuned enthusiastically and pretty. The tests affirm the success effect of the changes of heuristic elements. Moreover, the experiments show that the offered approach can make contributions to increase the gadget typical performance.

“Analysis of Optimal Thread Pool Size by Yibei Ling & Tracy Mullen”

This paper provides thread pool version on numerical announcement and provided a numerical shape for formulating the most useful pool duration for inexperienced thread supervision. Thread introduction time and the thread background switching time is used for measuring gold fashionable volume of the pool. But those parameters are hard to calculate at the intensity of kernel. Yet no internet server can calculate those parameters.


This paper recommend version fuzzing, a way that mixes size and model to make available a specific and measureable assessment of RM configurations. Model fuzzing starts by constructing a structure version that relates useful aid allocations to performance of the controlled device. For the net thread pool, that is articulated as a uni-modal concurrency throughput curve, a characteristic that is easy to approximate the usage of simple equations. Then, studies are carried out in which the RM execution engage with the system version to bet the performance of managed machine for RM configurations. Thru “fuzzing” the parameter of the gadget model, we explore the general performance of RM configurations for masses versions of the managed gadget and its workloads. This paper enlarge a technique for configuring RMs that makes use of version fuzzing and considers many evaluation criteria (e.g., excessive throughput, low quantity of threads).

“Prediction-Based Dynamic Thread Pool Management ofAgent Platform for Ubiquitous Computing by Ji Hoon Kim, Seungwok & Han, Hyun Ko”

The applications for ubiquitous device want to efficaciously using the belongings allocated within the environment and provide smart offerings to the users. As a manner to satisfy such necessity, the programs want to be evolved the use of intelligent marketers offering optimized offerings to all person. Moreover, the platform itself wishes that allows you to help platform degree service optimization. This paper proposed a prediction primarily based dynamic thread pool handling scheme using Gaussian distribution for
maximizing the usage of resources. The experiments shows that the proposed scheme outperforms the winning consultant thread pool fashions. It additionally added a modern-day shape of agent platform that gives interactive environment among smart dealers, at the side of the operation series flow of registration, elimination, and transmission of entrepreneurs with the proposed agent platform middle. It lets in every the institution management and man or woman control of the entrepreneurs for efficiency. Because the destiny paintings, we’re capable of similarly optimize the dynamic thread pool management scheme by using properly adjusting the parameters used in the prediction.

“Prediction-based Dynamic Thread Pool Scheme for Efficient Resource Usage by DongHyun Kang & Saeyoung Han”

This paper has progressed a few vulnerable factors of the static worker thread pool model and the requirement based totally thread version, in order that it could enthusiastically exchange the dimensions of thread swimming pools consistent with the purchaser’s requests. It is able to maintain a low reaction time to the requests, and is advanced to use the system sources proficiently. As an give up result, it predicts the wide variety of required threads earlier the usage of the performed exponential average scheme, and creates those threads earlier; in order that it improves the eliminate of the reaction time even as clients request, and it deletes the pointless threads for other packages to use the gadget property even as there are much less users’ requests. The ones performances are examined in the preceding phase using the cautioned scheme. Usual, our counseled version answers the proper architecture of device. While there are too many users’ requests, it extensively speaking attempts to lower the response time no matter the truth that it could use more resources. While there are handiest few requests, it returns all unused assets so they can be applied in extraordinary processes. In this paper, we test our version with a case with many requests and a case with few requests, and observe the end result of instances with watermark thread pool model. However, extra researches are desired to reveal that the prediction-based totally dynamic thread pool version is more efficient than the watermark model by way of manner of measuring the performance of the combination of times above for an extended time frame, or the overall performance of a well-known server surroundings which offers numerous services concurrently.

“A Novel Predictive and Self–Adaptive Dynamic Thread Pool Management by Kang-Lyul Lee, & Hee Yong Youn”

This paper added a unique technique which enthusiastically adapts the thread pool to the operation situation to maximize the device overall performance. To benefit the cause, the TEMA idea was proposed to as it should be expecting the quantity of threads. In addition, a brand new prediction primarily based thread pool employer changed into developed to dynamically alter the quantity of threads within the pool, thinking about the idle timeout duration and the conditions of thread destruction. It permits powerful thread advent & destruction consistent with the modern day repute. A test shows that the proposed scheme significantly outperforms the present schemes for numerous request fees in phrases of response time and CPU overhead. For destiny paintings, we plan to increase a brand new prediction scheme enhancing the prediction accurateness. Furthermore, scheme thinking about the priority of the watcher and employee thread is probably investigated to efficiently arbitrate the operations of
thread creation & destruction in conjunction with the call for processing.

“Design of Hierarchical Thread Pool Executor for DSM by Sitharamaiah Ramisetti & Rajeev Wankar”
This paper explains Hierarchical Thread Pool Executor by means of the non-blocking line has been considered and replicated at the cluster surroundings the usage of Min Heap tree. One of the essential obstacles of using the blockading queue is that the thread in multithreaded surroundings is blocked till it's far done inflicting performance deprivation. The anticipated design of Hierarchical Thread Pool Executor the use of the non-blocking queue overcomes this hassle and gives higher performance.

“Mobile Agent Based Elastic Executor Service by Anirban Bhattacharya”
This paper explains the idea to give reference structure to enforce a allotted executor provider the usage of JADE platform for cellular sellers. However it'd want extra designated tests with different mobile agent platform as properly. Tests also are necessary to be finished the way it reacts to failover in distinctive conditions. In destiny, mobile dealers in an executor service may be made wiser with changeable functionality. The nodes may be made to shuffle the request within the nodes relying on various capability of the node based. by way of assuming that each one containers/JVM are of same ability which can be exploited extra and solution may be advanced with different functionality of packing containers and extraordinary set of rules for capacity calculation.

“FREQUENCY BASED OPTIMIZATION STRATEGY FOR THREAD POOL SYSTEM by F. BAHADUR& M.NAEEM”
The primary contribution of this paper is the management of FBOS, a dynamic tuning approach it's primarily based completely on set of measureable procedures that are without complexity enormous and experimentally provable. FBOS approach is furnished for the ones server aspect packages that use thread pool structure. FBOS approach is applied in JAVA. This method dynamically resizes the thread pool at the idea of request frequency and it we ought to the thread pool device jogging gracefully. The quantitative measures of FBOS technique are turnaround time of requests, idle time of requests and the system in the course of. FBOS plays dynamic optimization of the thread pool when it famous that the turnaround time of jobs involve geared up time and it then react by growing the pool length in keeping with modern request frequency and then recycles the threads for arriving requests efficaciously.

“OPTIMIZING FREQUENCY-BASED THREAD POOL SYSTEM BY NON-BLOCKING QUEUES AND AUTOMATED TIMERS by Ghazala Ashraf, Faisal Bahadur, Mohammad Abrar Khan & Arif Iqbal Umar”
This paper introduces the automated timers to manage non-blocking thread pool system namely, non-blocking FBOS. Frequency-Based thread pool system (FBOS) has been redesigned by replacing Lock based data structures and objects with non-blocking ones. Non-blocking FBOS with automated timers has been designed by using spin primitives on non-blocking queues and atomic objects. We have observed the effects of non-blocking algorithms over the performance of existing scheme FBOS, by using non-blocking algorithms that use Compare and Swap (CAS) instructions in Java5. We have implemented frequency based automated timers instead of constant timers.

The results show that by using non-blocking queues / algorithms and automated timers in non-blocking FBOS with automated timers strategy, performance can be improved and better response time can be provided than
the one provided by FBOS strategy on high request arrival rate.

“Prediction and Frequency Based Dynamic Thread Pool SystemA Hybrid Model by Sumat Nazeer, Faisal Bahadur , Mohammad Abrar Khan, Abdul Hakeem, Miraj Gul, Arif Iqbal Umar”

This paper introduces a strategy which dynamically alters the size of a thread pool depending upon the rate of requests arrival, which results in the improvement of outcome and performance of the system. Hybrid (combination of prediction and frequency based) model is developed to gain the desired output. This model predicts the number of threads based on incoming frequencies.

This model also enthusiastically tunes the pool size by giving deliberation to the inactive time period. Threads are automatically destroyed due to increase in inactive time.

This methodology keeps away from synchronization overhead by considering the approaching solicitation frequencies rather than string pool estimate. By escaping synchronization overhead, its forecasts are significantly more precise than those of the TEMA technique. The subsequent exact forecast shuns the lacking of threads and prompts to execution change. The objective of this research was to enhance the reaction time of client’s requests, which has been effectively accomplished as reenactment results.

“Implementing Saturation Point Detection Mechanism in Frequency Based Thread Pool by Shahid Hussain”

This paper explains the implementation of Saturation point Detection Mechanism in FBOS. RFBOS is define and done with the aid of utilizing Java. This paper exhibited approach for servers that are running on multi threading. This plan is supposed for vertical scaling i.e. such as making ready asset a solitary server. It isn’t for horizontal scaling i.e. for corporations. It is define to be applied for the ones jobs which might be I/O sure or combination of I/O sure jobs & CPU sure jobs. This scheme cannot be applied for jobs which might be simply CPU bound jobs. Through implementing Saturation factor Detection Mechanism in Frequency based Thread gives brief reaction time to the customers through diminishing keep up time of each customer.

III. MOTIVATION

The motivation of our work is that Web servers, file servers and application servers implemented by FBOS can result in low performance when client’s requests are long running tasks i.e. service time of request is high, because FBOS is designed to optimize the pool when the requests are of low I/O intensity and those objects on the server that takes long time to service will slow down the system’s performance as FBOS only considers frequency of requests and not the service time of requests. Whenever a request of high I/O intensity i-e service time of the request is greater than one second than FBOS fails in that condition. This research is taken out to remove scheduling overhead from FBOS and to design and develop distributed framework for RFBOS

IV. PROBLEM SPECIFICATION

FBOS (frequency based optimization strategy) is single server based strategy, in which the context switching is performed. Second problem in FBOS is Synchronization overheads due to which system performance become slow.

V. PROPOSED SOLUTION

To implement the DFBOS we need a Master Server connected with multiple Slave Servers. These Servers are connected with each other by the help of Sockets as shown in Figure 5.1. Master server is used to process the Client’s requests by using the Thread Pool Strategy. In Master Serer thread pool size is restricted by the help of saturation point detection mechanism. When the pool
size reaches the saturation point, further requests of clients migrate to the 1st Slave Server and when the 1st Slave Server reaches its saturation point then further requests migrate to the 2nd Slave Server and so on.

**Figure 5.1 Network Design**

SDM (Saturation Detection Mechanism) is used to track saturation Point. Saturation point is detected by measuring the performance metrics on Master Server and Client Servers. The point at which maximum performance is obtained is marked as Saturation Point. RFM (Request Forwarding Mechanism) is used to migrate client’s requests from Master Server to the slave server. RFM uses Round Robin Scheduling to manage requests.

FOBS is a dynamic tuning approach to Optimize Pool length. FBOS strategy makes use of a framework that is developed in JAVA. This approach can dynamically resize the thread pool on the idea of request frequency. Global Table is used to save the IP address of the connected slave servers and also update the Saturation Point Status respectively.

**Figure 5.2 FBOS Server**

DFBOS system consists of one master FBOS server and more than 1 slave FBOS servers. First we will start Master server of DFBOS System as depicted in figure 5.2. When it is started it will start listening into two ports, the port2 is used to listen requests from slave servers. When we start any slave from network computer we will give the IP of Master server and slave will try to connect to the Master server on Port2, if it is a new slave then the Global Table is updated on Master Server to store the IP of slave as well as its saturation point (SP) is recorded as label “No”, i.e. Saturation point is not reached. When the saturation point of slave is reached it will send a signal to the Master server which will search its IP and update SP status to “YES” as well as the maximum capacity of slave. Port1 is used for the users/clients of DFBOS system that are sending requests. Master FBOS server can process the requests at specific frequency until its saturation point arrives, i.e. For example if request rate is 1000 requests per second and FBOS performance is not degraded then 1000 is not the saturation point and if request rate becomes 1200 request per second and throughput is degraded it means that 1200 is saturation point and server capacity is only 1000 requests per second. All other requests would be forwarded to Request Forwarding mechanism (RFM). Saturation Detection Mechanism will continuously check server throughput after every second and when it will detect low throughput than desired then it will mark the saturation point of FBOS server. FBOS will forward all those requests to RFM which are beyond its capacity. RFM component of FBOS server runs in round robin fashion i.e. it will pick a slave from Global table and forward one request to it and then it will pick next slave from table and forward it one request and so on. It will repeat this loop until all the requests beyond server capacity have sent to the salves. In this way it will
maintain the load balancing on the network. For example if request rate is 1200 requests per second and server capacity is 1000 then 200 requests would be send to RFM and if there are only two slaves connected to the server then each one will process 100 requests in every second. In this way all the slaves will be given equal load to process until a slave complains for its own saturation point. For example if request rate is 2000 r/s and there are 2 salves then each slave is responsible to process 500 requests every second, now if request rate suddenly increased from 2000 to 3000r/s then each salve will receive 1000 requests every second and if SDM of slave1 detects its SP at this rate then salve 1 will send signal to server about its SP and its capacity is also marked in the Global Table. In this case RFM will not send the requests to slave 1 more than its capacity.

VI. ANALYSIS AND RESULTS

In this section a comparison between DFBOS, RFBOS & FBO S is performed. Comparison was performed by keeping four aspects i.e Load Generation, Throughput, Response Time & Pool Size of the above motioned techniques.

Load Generation

On the below graph, Number of Requests are plotted at Y-Axis and Time in Seconds is plotted at X-Axis. We have Poisson appropriation having λ =25, 50, 75,100,125. We have the Workload of 100 milliseconds.

Throughput

On the below graph, Time is plotted at X Axis & Throughput is plotted at Y Axis. At the beginning, Throughput of all strategies is identical because of Non Restriction or Non Saturation Point detection. After attaining the SP i.e. throughput of 75 jobs / second. The throughput of RFBOS was constrained and the throughput of FBO S became reduced. While the throughput in DFBOS technique increases.

Figure 5.3 FBOS Slave Server

Architecture of slave server of FBOS is given in Figure 5.3. When this server starts it will connect to the master server and it will accept requests from the master server its SDM will repeatedly checks its maximum capacity( after every second ) on which it can process request with particular rate on high performance. When SDM detects it’s SP it will signal Master server that its capacity is finalized and Master will not send requests more than its capacity.
Average Throughput generated in FBOS is 45 requests per second. In RFBOs average Throughput were 60 requests per second while in DFBOS the average Throughput was 75 requests per second.

**Response Time**

In the below graph, Response / Jobs are plotted at X Axis whereas the Response time is plotted at Y Axis. By increasing the numbers of Jobs, the response time of FBOS and RFBOS is suffered while in DFBOS the response time is improved.

![Response Time Graph]

**Figure 6.3 Response Time Graph**

Average Response Time of FBOS technique is 125, whereas the average Response Time of RFBOs is 116 and average Response Time of DFBOS is 107. Hence the average response time of DFBOS is improved.

**Pool Size**

In the below graph, Time is plotted at X Axis and Pool size is plotted at Y Axis. Pool Size is measured in Number of Threads whereas Time is measured in Seconds. Pool Size is Restricted in RFBOs, due to which the Pool Size remain stable after reaching SP. In FBOS strategy Pool Size is not Restricted. Whereas in DFBOS pool size is also restricted, but in DFBOS tasks are migrated to next Server after reaching SP.

![Pool Size Graph]

**Figure 6.4 Pool Size Graph**

**Results Comparison**

The below table illustrate the comparison between FBOS, DFBOS and RFBOs.

<table>
<thead>
<tr>
<th></th>
<th>Response Time</th>
<th>Throughput</th>
</tr>
</thead>
<tbody>
<tr>
<td>DFBOS</td>
<td>14.4% more</td>
<td>DFBOS give 55% more throughput w.r.t FBOS.</td>
</tr>
<tr>
<td>DFBOS</td>
<td>7.7% more</td>
<td>DFBOS give 16% more throughput w.r.t FBOS.</td>
</tr>
</tbody>
</table>

**VII. CONCLUSION & FUTURE WORK**

The most important task of this research is to get better result of thread pool by developing distributed environment. The prior methods of Thread pool are single server and due to limited space, pool size was bounded. Due to restricted size the overhead threads goes into waiting state. DFBOS is developed to overcome these problems by using multiple servers i.e. One Master Server connected with Numbers of Client Servers. Master server detects the saturation point of each server and migrates the jobs to next server. To elaborate the results and assumption a JAVA framework called Thread Pool Tester is developed which is used as a simulation tool. Thread Pool Tester is used to test the performance. Thread pool tester also evaluates different throughput, frequencies, pool size and response time with FBOS and RFBOs and gives the better results in all aspects.

**Future Work**

In DFBOS method there is no procedure to share out the workload along with different servers. To uniformly
share out the request among the special nodes a Load Balancing Framework can be developed.

REFERENCES


